MPI Part 1

* References:

<u>Using MPI</u>. Gropp, Lusk Skjellum

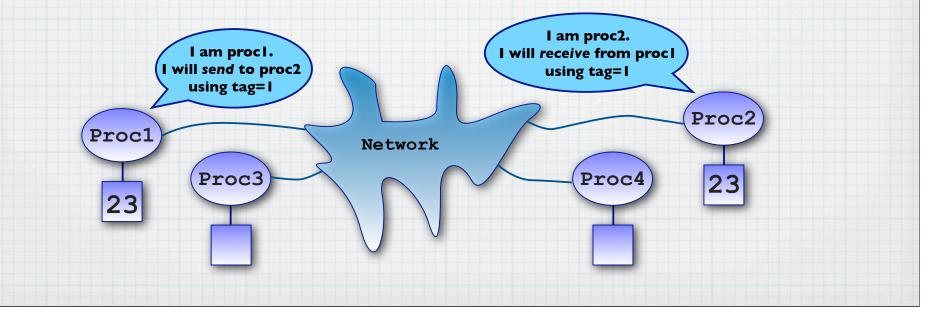
http://www.mpi-forum.org/docs/mpi-11-html/node182.html

* What is MPI?

- 1) MPI allows a collection of processes to communicate with messages.
- 2) MPI is a library of subroutines called from Fortran, C and C++. Programs are compiler with ordinary compilers and linked with the MPI library.
- 3) MPI is a specification which is independent from particular implementations. An MPI program should be portable to any vendors hardware that supports MPI.

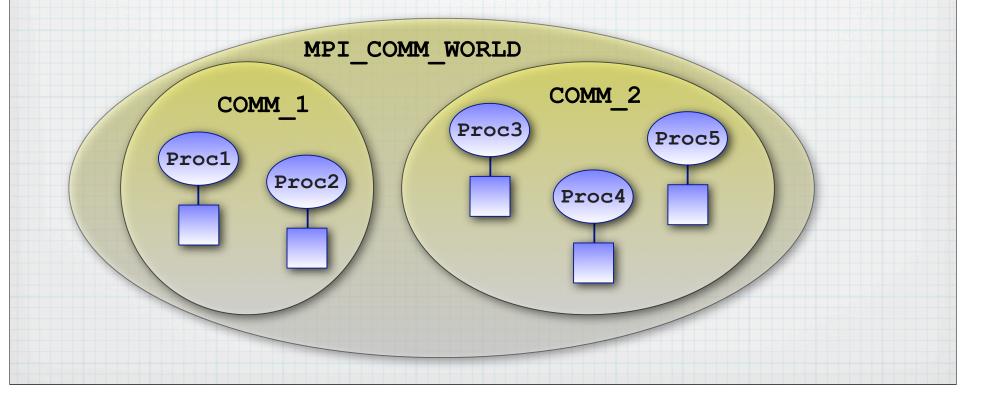
A minimal message-passing model

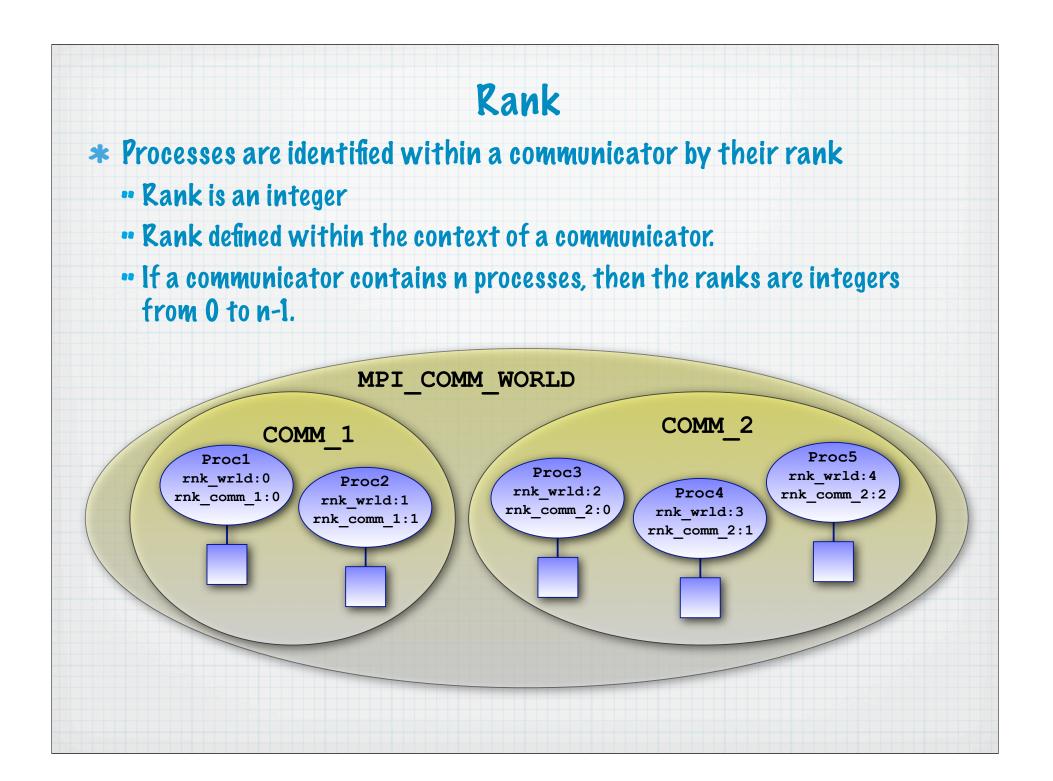
- * The processes execute in parallel and have separate address spaces.
- * Communication is cooperative. A message requires one process to execute a send command, and one process to execute a receive command.
- * Information from one process's address space (memory) is transfered to another address space (memory) using a message.
- * The two processes involved in the communication must agree upon a message tag to distinguish a message from other messages.



Communicators

- * Groups of processes are called communicators.
 - The default communicator is called MPI_COMM_WORLD. This communicator contains all the processes in the current MPI universe.
 - •• MPI allows for the formation of communicators within the global communicator.
 - •• Message tags are defined within the context of a communicator.





The "hello world" Program

- * Important features of the hello_world program
 - 1. Use the mpi module, or include the include file called mpif.h
 - 2. Initialize the MPI environment.
 - **3.** Determine how many processes are in the current MPI environment.
 - 4. Determine rank within the MPI_COMM_WORLD communicator
 - **5**. Terminate the MPI environment

Running the hello_world Program

PROGRAM hello_world USE mpi

CALL MPI_INIT (ierr)! initialize MPI environmentCALL MPI_COMM_SIZE (MPI_COMM_WORLD,npe_wrld,ierr)! determine world sizeCALL MPI_COMM_RANK (MPI_COMM_WORLD,rnk_wrld,ierr)! determine rank within world

PRINT "(A19,I3,A4,I4)"," hello from proc = ",rnk_wrld," of ",npe_wrld

```
CALL MPI FINALIZE (ierr)
```

! terminate MPI environment

END PROGRAM hello world

00	Terminal — csh — csh (ttyp1) — $80x14$	
bliss 39 ≻ mpiru	ın -np 8 hello_world	5
hello from prod	:= 2 of 8	
hello from prod	c = 1 of 8	
hello from prod	c = 0 of 8	
hello from prod	:= 5 of 8	
hello from prod	c = 6 of 8	
hello from prod	c = 3 of 8	
	c = 4 of 8	
hello from prod	c = 7 of 8	
bliss 40 >		
		A
		X
		1.

The slightly modified "hello world" Program

* Important features of the slightly modified hello_world program

- I. Use the mpi commands MPI_GET_PROCESSOR_NAME to determine where a processes is actually running.
- 2. Use the mpi commands MPI_WTICK and MPI_WTIME to time code
- **3.** Use the mpi commands MPI_BARRIER write output in order.

The slightly modified "hello world" Program

* Code (continued) for the slightly modified hello_world program

CALL MPI GET PROCESSOR NAME (proc name, name len, ierr)

```
wall_tick = MPI_WTICK () ! wall clock timer increment in seconds
IF (rnk_wrld == 0) PRINT "(A13,F12.8)"," wall_tick = ",wall_tick
```

```
! do some useless work
```

```
time_start = MPI_WTIME () ! wall clock timer start
```

```
x = 0.0_8
```

DO j = 1,5000

```
DO i = 1,5000
```

```
x = x + SIN (x+FLOAT (rnk wrld))
```

```
ENDDO
ENDDO
```

```
time end = MPI WTIME () ! wall clock timer stop
```

```
! write the results
```

```
DO n = 0, npe wrld-1
```

```
IF (rnk_wrld == n) THEN
```

```
PRINT "(A19, I3, A4, I4, A12, A16, A10, F8.5, A10, F12.8)", &
```

```
" hello from proc = ",rnk_wrld," of ",npe_wrld, &
```

```
" running on ", TRIM (proc_name), &
```

```
" time = ",time_end-time_start," answer = ",x
```

ENDIF

```
CALL MPI_BARRIER (MPI_COMM_WORLD,ierr)
```

ENDDO

CALL MPI FINALIZE (ierr)

END PROGRAM hello world 2

Collective Communication: Scatter

* Transfer information for one process to many (scatter) or collect information from many processes to one (gather)

* MPI_BCAST broadcasts a message from the process with rank ROOT to all processes of the communicator group COMM, itself included. It is called by all members of group using the same arguments. On return, the contents of root's send buffer has been copied to the receive buffer on all processes.

MPI_BCAST (buffer,data_count,data_type,root,comm)

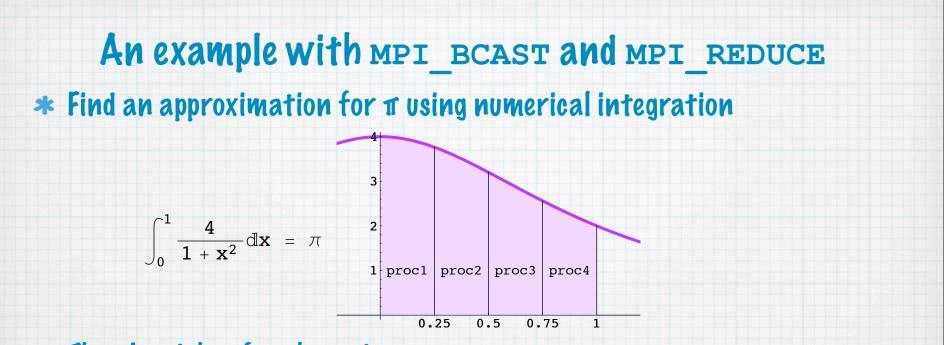
Collective Communication: Gather

* MPI_GATHER: Each process (root process included) sends the contents of its send buffer to the root process. The root process receives the messages into the receive buffer and stores them in rank order.

MPI_GATHER(send_buffer,send_count,send_type, recv_buffer,recv_count,recv_type,root,comm,ierr)

* MPI_REDUCE: Combines the elements in the send buffer of each process in the communicator group comm, using the operation op, and returns the combined value in the receive buffer of the process with rank root.

where op can be several things including MPI_MAX (MPI_MIN) for maximum (minimum), MPI SUM for summation.



* The algorithm for the code:

- 1. The root process will read the global number of intervals and broadcast the number to the other processes using MPI_BCAST.
- 2. Each process will then determine its subinterval using its rank in the communicator and integrate to find its subarea
- **3.** Using MPI_REDUCE with the option MPI_SUM the subareas are summed to find the total area

	Code the Pi example
	PROGRAM pi
	USE mpi
	IMPLICIT NONE
	INTEGER :: npe_wrld,rnk_wrld,n,i,ierr
	REAL (KIND=SELECTED_REAL_KIND (12)) :: &
	del_x,x_left,pi_piece,pi_approx,time_start,time_end,x
! :	setup MPI
	CALL MPI_INIT (ierr)
	CALL MPI_COMM_SIZE (MPI_COMM_WORLD, npe_wrld, ierr)
	CALL MPI_COMM_RANK (MPI_COMM_WORLD, rnk_wrld, ierr)
! :	read and broadcast total number of intervals
	IF (rnk_wrld==0) THEN
	PRINT *, 'Enter the total number of intervals '
	READ (*,*) n
	ENDIF
	CALL MPI_BCAST (n,1,MPI_INTEGER,0,MPI_COMM_WORLD,ierr)
	<pre>time_start = MPI_WTIME () ! wall clock timer start</pre>
! :	integrate subinterval
	del_x = 1.0_8/DBLE (n); x_left = DBLE (rnk_wrld)/DBLE (npe_wrld);
	pi_piece = 0.0_8
	DO i = 1,n/npe_wrld
	$x = x_{left} + del_x * (DBLE(i) - 0.5_8)$
	pi_piece = pi_piece + del_x*(4.0_8/(1.0_8 + x**2))
	ENDDO
! (gather the pieces of the pi
	CALL MPI_REDUCE(pi_piece,pi_approx,1,MPI_DOUBLE_PRECISION,MPI_SUM,0, &
	MPI_COMM_WORLD,ierr)
	<pre>time_end = MPI_WTIME () ! wall clock timer stop</pre>
!]	print the approximate value
	IF (rnk_wrld==0) THEN
	PRINT "(A12,F22.20)","pi_approx = ",pi_approx ENDIF
	CALL MPI BARRIER (MPI COMM WORLD, ierr)
	PRINT "(A12,F14.10)","time = ",time_end-time_start
	CALL MPI FINALIZE (ierr)

Point-to-point communication

* Here we send messages directly form one process to another.

* MPI_SEND:

- This is a blocking send. Control does not return until the message data has been safely stored away so that the sender is free to overwrite the send buffer.
- •• The syntax of the blocking send operation is given below:

MPI_SEND (BUFFER, DATA_COUNT, DATA_TYPE, DEST, TAG, COMM, IERR)

where

DEST is the rank of destination (integer) within COMM TAG is the message tag (integer)

Point-to-point communication

* MPI_RECV:

- This is a blocking receive. Control returns only after the receive buffer contains the newly received message.
- •• The syntax of the blocking send operation is given below:
- MPI_RECV (BUFFER, DATA_COUNT, DATA_TYPE, SOUR, TAG, COMM, STATUS, IERR)

where

SOUR is the rank of source (integer) within COMM. The source can also be specified as MPI_ANY_SOURCE TAG is the message tag (integer). The tag can also be specified as

MPI ANY TAG

*This is a "master-slave" algorithm. One process (the master) is responsible for the coordinating the work of the others (the slaves).

* We wish to perform a matrix-vector multiply in parallel.

 $\mathbf{A}b = c$

* The <u>masters</u> algorithm for the code:

1. The master will broadcast the vector **b** to all the slaves.

2. The master will send one row of the matrix A to each slave.

3. The master then waits for the slave to perform the dot product and return the element of c. At this time the master sends that slave a new row of A. Continue until all rows are processed.

- * The <u>slaves</u> algorithm for the code:
 - 1. The slaves receive vector b from master.

2. Perform dot-products of b and rows of A. Send result to master

* The code is clearly partitioned into a master part and a slave part

```
PROGRAM mat vec
USE mpi
IMPLICIT NONE
INTEGER, PARAMETER :: rows=100, cols=100
INTEGER :: npe wrld, rnk wrld, master, i, j, count rows, sender, row index, ierr
INTEGER :: status (MPI STATUS SIZE)
REAL (KIND=SELECTED REAL KIND (12)) :: &
   a(rows,cols),b(cols),c(rows),buffer(cols),ans,time start,time end
CALL MPI INIT (ierr)
CALL MPI COMM SIZE (MPI COMM WORLD, npe wrld, ierr)
CALL MPI COMM RANK (MPI COMM WORLD, rnk wrld, ierr)
master = 0
IF (rnk wrld==master) THEN ! THE MASTER DOES THIS BLOCK OF CODE
ELSE ! THE SLAVES DO THIS BLOCK OF CODE
  •
ENDIF
CALL MPI FINALIZE (ierr)
END PROGRAM mat vec
```

Point-to-point communication. Matrix-vector multiplication. * The master code looks like this:

```
DO j = 1, cols ! make an arbitrary matrix a and vector b
   b(i) = 1.08
   DO i = 1, rows
      a(i,j) = DBLE(i+j)
   ENDDO
ENDDO
CALL MPI BCAST (b,cols,MPI DOUBLE PRECISION,master,MPI COMM WORLD,ierr)
count rows = 0
DO i = 1, npe wrld-1
   DO_j = 1, cols
      buffer(j) = a(i,j)
   ENDDO
   CALL MPI SEND (buffer, cols, MPI DOUBLE PRECISION, i, i, MPI COMM WORLD, ierr)
   count rows = count rows+1
ENDDO
DO i = 1, rows
   CALL MPI RECV (ans,1, MPI DOUBLE PRECISION, &
                     MPI ANY SOURCE, MPI ANY TAG, MPI COMM WORLD, status, ierr)
   sender
             = status (MPI SOURCE)
   row index = status (MPI TAG) ! tag value in status is the row index
   c(row index) = ans
   IF (count rows < rows) THEN ! more work to be done. send another row
      DO_j = 1, cols
         buffer(j) = a(count rows+1,j)
      ENDDO
      CALL MPI SEND (buffer, cols, MPI DOUBLE PRECISION, &
                                    sender,count rows+1,MPI COMM WORLD,ierr)
      count rows = count rows+1
   ELSE ! tell sender that there is no more work
      CALL MPI SEND (MPI BOTTOM,0, MPI DOUBLE PRECISION, sender,0, MPI COMM WORLD, ierr)
   ENDIF
ENDDO
```

* The slave code looks like this:

```
CALL MPI BCAST (b, cols, MPI DOUBLE PRECISION, master, MPI COMM WORLD, ierr)
DO
   CALL MPI RECV (buffer, cols, MPI DOUBLE PRECISION, master, &
                                     MPI ANY TAG, MPI COMM WORLD, status, ierr)
   IF (status(MPI TAG)==0) EXIT ! there is no more work
   row index = status (MPI TAG) ! tag value status is the row index
   ans = 0.08
   DO i = 1, cols
      ans = ans + buffer(i)*b(i)
   ENDDO
   CALL MPI SEND (ans,1, MPI DOUBLE PRECISION, &
                                       master,row index,MPI COMM WORLD,ierr)
ENDDO
```

* Running the code

* slower with more processes...

time = 0.0922520161 time = 0.0922918320 eden 11 > mpirun -np 3 mat_v time = 0.2493729591 time = 0.2492809296 time = 0.2493259907 eden 12 > mpirun -np 4 mat_v time = 0.5544278622 time = 0.5542669296			
time = 0.2493729591 time = 0.2492809296 time = 0.2493259907 eden 12 > mpirun -np 4 mat_y time = 0.5544278622			
time = 0.2492809296 time = 0.2493259907 eden 12 > mpirun -np 4 mat_v time = 0.5544278622	ec		
time = 0.2493259907 eden 12 > mpirun -np 4 mat_v time = 0.5544278622	ec		
eden 12 > mpirun -np 4 mat_v time = 0.5544278622	ec		
time = 0.5543239117			
time = 0.5542860031			
eden 13 >			
			X
			T

	Next time	2	
1. Non-blocking sends a	nd receives		
•• Overlapping comm	unications and com	putations	
2. Topologies			
3. MPI datatypes			